

# On families of weakly cross-intersecting set-pairs

Zoltán Király<sup>a,1</sup>, Zoltán L. Nagy<sup>a</sup>, Dömötör Pálvölgyi<sup>a</sup>, Mirkó Visontai<sup>b,2</sup>

<sup>a</sup>*Department of Computer Science, Eötvös University Budapest, Pázmány Péter sétány 1/C. H-1117 Hungary*

<sup>b</sup>*Department of Mathematics, University of Pennsylvania, 209 S. 33rd Street, Philadelphia PA, 19104 USA*

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## 1. Introduction

Let  $\mathcal{F}$  be a family of pairs of sets. We call it an  $(a, b)$ -set system if for every set-pair  $(A, B)$  in  $\mathcal{F}$  we have that  $|A| = a$ ,  $|B| = b$ ,  $A \cap B = \emptyset$ . The following classical result on families of cross-intersecting set-pairs is due to Bollobás [6]. Let  $\mathcal{F}$  be an  $(a, b)$ -set system with the cross-intersecting property, i.e., for  $(A_i, B_i), (A_j, B_j) \in \mathcal{F}$  with  $i \neq j$  we have that both  $A_i \cap B_j$  and  $A_j \cap B_i$  are non-empty. The maximum possible size of such a set system is  $\binom{a+b}{a}$ , independent of the size of the ground set, and this bound is sharp. Surprisingly, the same upper bound holds even if we relax the cross-intersecting property, namely if for  $(A_i, B_i), (A_j, B_j) \in \mathcal{F}$  we only require that  $A_i \cap B_j \neq \emptyset$  when  $i < j$ , as was shown in [7]. Several further generalizations were investigated in [2, 8, 10, 12]. For more details on the history and applications of this problem we refer to the surveys [14, 15] and Chapter 1 of [1].

In this paper we consider the following variant of the problem, introduced by Tuza [13].

**Definition 1.** *Let  $\mathcal{F}$  be an  $(a, b)$ -set system.  $\mathcal{F}$  is weakly cross-intersecting if for any  $(A_i, B_i), (A_j, B_j) \in \mathcal{F}$  with  $i \neq j$  we have that  $A_i \cap B_j$  and  $A_j \cap B_i$  are not both empty.*

We investigate the maximum possible size of such a system, which we denote by  $g(a, b)$ . In Section 2 we give an explicit construction based on lattice paths and prove our main result that  $\liminf_{a+b \rightarrow \infty} g(a, b) / \binom{a+b}{a} \geq 2 - o(1)$ . Here, and throughout the paper, we use  $o(1)$  to denote that an expression tends to zero as  $a+b$  tends to infinity. This is the first construction giving this lower bound, which is stronger than the previously best known  $\left(2 - \frac{ab}{(a+b)(a+b-1)}\right)$ . In Section 3 we recall the known upper bounds and introduce a fractional version of the problem. We prove an upper bound for the fractional version that matches the best known bound for the original problem. Finally, in Section 4 we present some computational results

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*Email addresses:* kiraly@cs.elte.hu (Zoltán Király), nagyzottanlorant@gmail.com (Zoltán L. Nagy), dom@cs.elte.hu (Dömötör Pálvölgyi), mirko@math.upenn.edu (Mirkó Visontai)

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for small values of  $a$  and  $b$ , in particular, we show that  $g(2, 2) = 10$ . We conclude by posing several interesting open problems.

Before we start, we mention some related problems that arise in Graph Theory.

### 1.1. Economical bipartite edge-coverings

Let  $K_N$  denote the complete graph on  $N$  vertices. Suppose we are given  $K_N$  and a positive integer  $k$ . We would like to cover the graph's edges with a collection of complete bipartite graphs,  $\mathcal{G}_j = \{(X_j \cup Y_j, E_j)\}_{j=1}^m$ , such that every edge of  $K_N$  is covered by at most  $k$  bipartite graphs. This problem was studied by Alon [3] and he showed that the minimum possible number of bipartite graphs to cover all the edges of  $K_N$  is  $m = \Theta(kN^{\frac{1}{k}})$ .

A related question is when we require every vertex – instead of every edge – to be covered by only a limited number of bipartite graphs. More precisely, suppose we are given  $K_N$  and two positive integers  $a$  and  $b$ . We would like to cover  $K_N$ 's edges with a collection of complete bipartite graphs,  $\mathcal{G}_j = \{(X_j \cup Y_j, E_j)\}_{j=1}^m$ , such that for every vertex  $v \in K_N$  there are at most  $a$  bipartite graphs such that  $v$  belongs to  $X_j$ , and at most  $b$  bipartite graphs such that  $v$  belongs to  $Y_j$ . Note that such a covering might not exist, if  $N$  is large. A natural question is to determine the largest  $N$  for which we still have a covering. The following proposition exposes the connection to the maximal size cross-intersecting families.

**Proposition 2.** *If  $\mathcal{G}_j = \{(X_j \cup Y_j, E_j)\}_{j=1}^m$  is a collection of complete bipartite graphs, that covers the edges of  $K_N$  in a way that every vertex is contained in at most  $a$  of the  $X_j$ 's and in at most  $b$  of the  $Y_j$ 's, then  $N \leq g(a, b)$ . Furthermore, this bound is sharp.*

*Proof.* We show how to construct a weakly cross-intersecting  $(a, b)$ -set system from such a covering. Let the elements of the ground set  $\{p_j\}_{j=1}^m$  correspond to the bipartite graphs in the covering. And let the pairs of sets  $\{A_i, B_i\}_{i=1}^N$  correspond to the vertices of the complete graph. Set  $p_j \in A_i$  if  $v_i \in X_j$  and set  $p_j \in B_i$  if  $v_i \in Y_j$ . This way we ensured that all the sets  $A_i$  have cardinality at most  $a$ , and the sets  $B_i$  have cardinality at most  $b$ . Also, any two pairs of sets weakly cross-intersect at the element corresponding to the bipartite graph that covered the edge connecting their corresponding vertices. Finally, adding some additional arbitrary points to the sets containing less than  $a$  (respectively,  $b$ ) elements, we get a weakly cross-intersecting  $(a, b)$ -set system.

Similarly, we can also construct a covering from a weakly cross-intersecting  $(a, b)$ -set system. This shows that the two problems are equivalent, so the upper bound is sharp.  $\square$

## 2. Lower bounds

We start with the following two claims made by Tuza [13].

**Claim 3.**  $g(a, 1) \geq 2a + 1$ .

*Proof.* Let  $B_i = \{i\}$  and  $A_i = \{i + j \pmod{(2a + 1)} \mid 1 \leq j \leq a\}$ , for  $i = 0, \dots, 2a$ .  $\square$

**Claim 4.**  $g(a, b) \geq g(a - 1, b) + g(a, b - 1)$ .

*Proof.* Suppose we have a construction  $\{(A'_i, B_i)\}$  of cardinality  $g(a-1, b)$  and another construction  $\{(A_j, B'_j)\}$  of cardinality  $g(a, b-1)$ . Let  $x$  be an element not contained in any of these sets. Then,  $\{(A'_i \cup \{x\}, B_i)\} \cup \{(A_j, B'_j \cup \{x\})\}$  gives a construction for a  $(a, b)$ -set system of cardinality  $g(a-1, b) + g(a, b-1)$ .  $\square$

From these claims one can easily obtain the following corollary.

**Corollary 5.**  $g(a, b) \geq 2\binom{a+b}{a} - \binom{a+b-2}{a-1} = \left(2 - \frac{ab}{(a+b)(a+b-1)}\right) \binom{a+b}{a}$ .

*Proof.* We prove by induction. Let  $\tilde{g}(a, b) = 2\binom{a+b}{a} - \binom{a+b-2}{a-1}$ . From Claim 3 we have that  $g(a, 1) \geq \tilde{g}(a, 1)$  for all  $a$ , and by symmetry that  $g(1, b) \geq \tilde{g}(1, b)$  for all  $b$ , which settles the base case. The inductive step follows from Claim 4, and the fact that  $\tilde{g}(a, b) = \tilde{g}(a-1, b) + \tilde{g}(a, b-1)$ .  $\square$

From this, one immediately gets that  $\lim_{a \rightarrow \infty} \tilde{g}(a, a) / \binom{2a}{a} = \frac{7}{4}$ , and hence an asymptotic lower bound on the ratio  $g(a, a) / \binom{2a}{a}$ . The following simple construction shows that the lower bound can be improved to  $\liminf_{a \rightarrow \infty} g(a, a) / \binom{2a}{a} \geq 2 - o(1)$ . Note that we use  $\liminf$  here as the limit might not exist.

**Claim 6.**  $g(a, a) \geq \left(2 - \frac{1}{a+1}\right) \binom{2a}{a}$ .

*Proof.* Consider a ground set of  $2a+1$  elements. We form a bipartite graph  $(X \cup Y, E)$  the following way. Each vertex in  $X$  corresponds to an  $a$ -subset of the ground set, and  $Y$  is just a copy of  $X$ . We connect vertices  $x \in X$  and  $y \in Y$  if the corresponding sets are disjoint. The resulting graph is regular, so, by König's Theorem, it must have a perfect matching  $\mathcal{M}$ . For every edge  $uv$  of  $\mathcal{M}$  we add a set-pair  $(A_i, B_i)$  to our system, where  $A_i$  corresponds to  $u$  in  $X$  and  $B_i$  corresponds to  $v$  in  $Y$ . It is easy to see that this family will be indeed weakly cross-intersecting, and has  $\binom{2a+1}{a} = \frac{2a+1}{a+1} \binom{2a}{a} = \left(2 - \frac{1}{a+1}\right) \binom{2a}{a}$  elements.  $\square$

In what follows, we give another explicit construction using lattice paths that improves on this bound and also works for general  $a$  and  $b$ . We obtain that  $\liminf_{a+b \rightarrow \infty} g(a, b) / \binom{a+b}{b}$  is also at least  $2 - o(1)$ .

### 2.1. A construction using lattice paths

Now we describe the construction of an  $(a, b)$ -set system  $\mathcal{F}$  for arbitrary  $a$  and  $b$  using lattice paths. Let  $m = 2a + 2b - 1$  be the cardinality of the ground set. We will construct  $\mathcal{F} = \{(A, B) : |A| = a, |B| = b, A, B \subseteq \{0, \dots, m-1\}\}$ .

Let  $\mathcal{L}(a, b)$  denote a set of all lattice paths on an  $a \times b$  grid from  $(0, 0)$  to  $(a, b)$ , where each move is either to the right, or up; and the path is strictly below the diagonal except at the two endpoints. We will identify a path  $\pi \in \mathcal{L}(a, b)$  with its binary representation, i.e., a permutation  $\pi$  of the multiset  $\{0^a, 1^b\}$ , using 0 and 1 for steps to the right and up, respectively. For each path  $\pi \in \mathcal{L}(a, b)$  we create the following  $m$  set-pairs  $\{(A_{\pi, i}, B_{\pi, i})\}_{i=0}^{m-1}$ . Let  $A_{\pi, i} = \{i+t \bmod m \mid 1 \leq t \leq a, \pi(t) = 0\}$  and  $B_{\pi, i} = \{i+t \bmod m, \mid 1 \leq t \leq a+b, \pi(t) = 1\}$ .

We have to show that the obtained  $\mathcal{F}$  indeed satisfies Definition 1. Clearly, it is an  $(a, b)$ -set system by construction. To show the weakly cross-intersecting property, first observe that  $A_{\pi,i} \cap B_{\sigma,i} \neq \emptyset$  for any  $\pi \neq \sigma$  as the paths are different and hence there exists a  $t$  such that  $\pi(t) = 0$  and  $\sigma(t) = 1$ . If  $i \neq j$  then  $A_{\pi,i} \cap B_{\sigma,j} \neq \emptyset$  or  $A_{\sigma,j} \cap B_{\pi,i} \neq \emptyset$  follows from the fact that any proper suffix of a lattice path in  $\mathcal{L}(a, b)$  cannot equal the prefix of another (or even the same) lattice path, and hence there is a  $t$  such that  $\pi(t) \neq \sigma(t + i - j)$ . One way to see this is to consider the ratio of the right and up steps: in the prefix this ratio has to be greater than  $a/b$  whereas in the suffix it has to be strictly less than that.

We note that for  $\gcd(a, b) = 1$  this  $(a, b)$ -set system is *maximal*, in the sense that we cannot add another pair of sets to it without violating the weakly cross-intersecting property.

Unfortunately, there is no simple closed formula for  $|\mathcal{L}(a, b)|$  for general  $(a, b)$ . There is a generating function for the number of lattice paths that never cross the diagonal but are allowed to touch it [11]. We believe, however, that it is unlikely that one can obtain better asymptotics using that result. For our purposes, the following result for the special case of relative primes  $a$  and  $b$  will suffice.

**Theorem 7** (Bizley [5]).  $|\mathcal{L}(a, b)| = \binom{a+b}{a}/(a+b)$  if  $\gcd(a, b) = 1$ .

From Theorem 7 and the construction described above we obtain a better lower bound than in Corollary 5.

**Theorem 8.**  $g(a, b) \geq (2 - \frac{1}{a+b})\binom{a+b}{a}$ , if  $\gcd(a, b) = 1$ .

*Proof.*  $g(a, b) \geq (2a + 2b - 1)|\mathcal{L}(a, b)| = (2a + 2b - 1)\binom{a+b}{a}/(a+b)$ . □

**Corollary 9.**  $g(a, a-1) \geq (2 - \frac{1}{2a-1}) \cdot \binom{2a-1}{a}$ .

Using this we can improve Claim 6.

**Corollary 10.**  $g(a, a) \geq (2 - \frac{1}{2a-1}) \cdot \binom{2a}{a}$ .

*Proof.* Claim 4 and Corollary 9 gives  $g(a, a) \geq 2g(a, a-1) = (2 - \frac{1}{2a-1}) \cdot \binom{2a}{a}$ . □

We can extend this idea to a general  $(a, b)$  by using Claim 4 repeatedly until we decrease  $a + b$  to a prime  $p$  for which we automatically have  $\gcd(p - q, q) = 1$  for  $0 < q < p$  and then invoke Theorem 8.

**Theorem 11.**  $\liminf_{a+b \rightarrow \infty} g(a, b) \geq (2 - o(1))\binom{a+b}{a}$ .

*Proof.* We start with a well-known result in Number Theory. There exists  $\delta > 0$  and  $x_0$  such that for all  $x > x_0$  the interval  $[x - x^{1-\delta}, x]$  contains a prime [9]. We will use this result for  $x = a + b$ . If either  $a$  or  $b$  is smaller than  $x^{1-\delta}$  then by Corollary 5 we have  $g(a, b)/\binom{a+b}{a} \geq 2 - \frac{x \cdot x^{1-\delta}}{x(x-1)} = 2 - o(1)$ . Otherwise, we will use the result that there is a prime  $p$  such that  $0 \leq x - p \leq x^{1-\delta} \leq a, b$ . By applying Claim 4 several times, and then Theorem 8, we obtain that

$$\begin{aligned}
g(a, b) &\geq \sum_{i=0}^{x-p} \binom{x-p}{i} g(a - (x-p) + i, b - i) \\
&= \sum_{i=0}^{a+b-p} \binom{a+b-p}{i} g(p - (b-i), b-i) \\
&\geq \sum_{i=0}^{a+b-p} \binom{a+b-p}{i} (2p-1) \binom{p}{b-i} / p \\
&= \frac{2p-1}{p} \binom{a+b}{b}.
\end{aligned}$$

And since  $(2p-1)/p \geq 2 - 1/(x - x^{1-\delta}) = 2 - o(1)$ , we have proven the theorem.  $\square$

### 3. Upper bounds

For the sake of completeness we also mention the upper bounds which appeared in [13].

**Claim 12.**  $g(a, 1) = 2a + 1$ .

For general  $a$  and  $b$  we do not have matching upper bounds. In fact, the following theorem (also appeared as an exercise in [4, p. 12]) gives the best known upper bound.

**Theorem 13** (Tuza, [13]).  $g(a, b) < \frac{(a+b)^{a+b}}{a^a b^b}$ .

*Proof.* The proof is a standard application of the probabilistic method. Let  $\{(A_i, B_i)\}_{i=1}^N$  be a weakly cross-intersecting  $(a, b)$ -set system. Take a random partition of all the elements of the ground set. For every element  $v$ , with probability  $\frac{a}{a+b}$  put  $v$  in  $X$  and with probability  $\frac{b}{a+b}$  put it in  $Y$ . We say that the set-pair  $(A_i, B_i)$  is *contained* in  $(X, Y)$ , if  $A_i \subseteq X$  and  $B_i \subseteq Y$ . The probability that a set-pair is contained in  $(X, Y)$  is  $\frac{a^a b^b}{(a+b)^{a+b}}$ . On the other hand, at most one set-pair  $(A_i, B_i)$  can be contained in any partition  $(X, Y)$ . Thus  $N \leq \frac{(a+b)^{a+b}}{a^a b^b}$ . Moreover, since not all random partitions can contain an  $(a, b)$  set-pair (e.g., when  $X = \emptyset$ ), we have a strict inequality.  $\square$

We can strengthen the above bound in some sense. Let  $g(a, b; m)$  denote the maximum possible size of a weakly cross-intersecting  $(a, b)$ -set system on a ground set of size  $m$ , and let  $k = k(a, b; m) := \lceil \frac{am-b}{a+b} \rceil$ . Then concentrating the distribution to sets  $X$  of size  $k$  we obtain the following theorem.

**Theorem 14.**  $g(a, b; m) \leq \frac{\binom{m}{k}}{\binom{m-a-b}{k-a}}$ .

As  $\frac{\binom{m}{k}}{\binom{m-a-b}{k-a}} < \frac{(a+b)^{a+b}}{a^a b^b}$  for  $m \geq a+b$ , we get back Theorem 13 as a corollary of this bound.

A reason why Theorem 13 is not that simple to improve on in general is, that this bound is optimal if we allow *fractional* weakly cross-intersecting families.

### 3.1. Fractional weakly cross-intersecting families

We now define the fractional relaxation of the weakly cross-intersecting property.

**Definition 15.** A weighted  $(a, b)$ -set system is an  $(a, b)$ -set system  $\{(A_i, B_i)\}_{i=1}^N$  with non-negative weights  $\lambda_i$  associated to each set-pair  $(A_i, B_i)$ .

**Definition 16.** An  $(a, b)$ -set system  $\{(A_i, B_i)\}_{i \in I}$  is separable if  $(\bigcup_{i \in I} A_i) \cap (\bigcup_{j \in I} B_j) = \emptyset$ .

**Definition 17.** A weighted  $(a, b)$ -set system  $\{(A_i, B_i)\}_{i=1}^N$  with weights  $\lambda_i$  is fractionally weakly cross-intersecting if for every  $I \subseteq \{1, \dots, N\}$  whenever  $\{(A_i, B_i)\}_{i \in I}$  is separable, then  $\sum_{i \in I} \lambda_i \leq 1$ .

We denote the maximum total value  $\sum_{i=1}^N \lambda_i$  of such a system on a ground set of size  $m$  by  $g^*(a, b; m)$ . Note that  $\lambda_i \leq 1$ , otherwise  $I = \{i\}$  would already violate the definition. If we require each  $\lambda_i$  to be an integer, then we get back Definition 1.

The following theorem shows that the bound of Theorem 14 is in fact optimal for the fractional case.

**Theorem 18.** If  $k = k(a, b; m)$  then  $g^*(a, b; m) = \frac{\binom{m}{k}}{\binom{m-a-b}{k-a}}$ .

*Proof.* First we show that this is an upper bound on  $g^*(a, b; m)$  as well. For a subset  $X$  with  $|X| = \binom{m}{k}$ , let  $I_X = \{j \mid A_j \subseteq X, B_j \cap X = \emptyset\}$ . By definition,  $\sum_{i \in I_X} \lambda_i \leq 1$  for all such  $X$ . A fixed  $j$  is in  $I_X$  for exactly  $\binom{m-a-b}{k-a}$  different choices of  $X$ .

For a construction, let the weighted  $(a, b)$ -set system consist of all possible disjoint pairs  $\{(A_i, B_i)\}_{i=1}^N$ , where  $|A_i| = a$ ,  $|B_i| = b$ , so we will have  $N = \binom{m}{a} \binom{m-a}{b}$ . For all possible values of  $i$ , let  $\lambda_i = \lambda = \frac{\binom{m}{k}}{\binom{m-a-b}{k-a} \binom{m}{a} \binom{m-a}{b}}$ . Let  $I \subseteq \{1, \dots, N\}$  be such that the sub-system  $\{(A_i, B_i)\}_{i \in I}$  is separable. We claim that  $\sum_{i \in I} \lambda_i \leq 1$ , i.e., that  $|I| \leq \frac{\binom{m-a-b}{k-a} \binom{m}{a} \binom{m-a}{b}}{\binom{m}{k}}$ . Let  $X = \bigcup_{i \in I} A_i$  and  $\ell = |X|$ . By definition,  $|I| \leq \binom{\ell}{a} \binom{m-\ell}{b}$ . An easy calculation shows that  $\binom{\ell}{a} \binom{m-\ell}{b}$  takes its maximum at  $\ell = k$ , and clearly  $\binom{m}{k} \binom{k}{a} \binom{m-k}{b} = \binom{m}{a} \binom{m-a}{b} \binom{m-a-b}{k-a}$ .  $\square$

## 4. Proving bounds using computer programs

Next we present a straightforward method for computing  $g(a, b)$  for fixed  $a$  and  $b$  with the help of a computer program. If we know  $m$ , the size of the ground set (or at least an upper bound on the size), then we can find a maximum size weakly cross-intersecting  $(a, b)$ -set system as follows. First, generate all possible  $(a, b)$  set-pairs. Next create a graph whose vertices represent set pairs and edges are drawn between vertices if the corresponding set-pairs weakly cross-intersect. Finally, find a maximum clique in this graph (since cliques in the graph correspond to set systems with weakly cross-intersecting property by construction).

Let us denote the size of the smallest ground set on which it is possible to realize a weakly cross-intersecting  $(a, b)$ -set system with  $N$  set-pairs by  $h(a, b; N)$  (if it exists). We are only interested in the case where  $N > \binom{a+b}{b}$ , consequently,  $h(a, b; N) > a + b$ .

**Claim 19.**  $\binom{h(a,b;N)}{2} \leq ab \cdot N$ .

*Proof.* Suppose we have a construction of size  $N$  on a set of size  $m$  and  $\binom{m}{2} > ab \cdot N$ . Define a graph on the ground set in the following way. Connect  $x$  and  $y$  if there is an  $i$  such that  $A_i$  contains one of them and  $B_i$  the other one. By our assumption, there must be a pair of points that are not connected to each other. Contracting these points gives a smaller ground set, where still any two set-pairs weakly cross-intersect. It might happen that  $x$  and  $y$  are both in some  $A_j$  (or  $B_j$ ), in this case just add to the shrunken set an arbitrary element that is not in  $A_j \cup B_j$ . Note that since  $m > a + b$  we can always choose such an element.  $\square$

**Theorem 20.**  $g(2, 2) = 10$ .

*Proof.* From Claim 4 we have that  $g(2, 2) \geq 10$ . So, now we have left to prove that  $g(2, 2) \leq 10$ . We prove by contradiction. Suppose, that there is a weakly cross-intersecting  $(2, 2)$ -set system with  $N = 11$  set-pairs. Then Claim 19 implies that this  $(2, 2)$ -set system has a realization on a ground set with at most 9 elements. We performed an exhaustive search on the ground sets of sizes up to 9 with our computer program and found that no such  $(2, 2)$ -set system has cardinality 11 or more. Hence, we can conclude that the maximum size of a  $(2, 2)$ -set system is 10.  $\square$

Unfortunately, the running time of this brute force search grows superexponentially and we could not compute further values. We found a  $(3, 2)$ -set system of size 19, hence  $g(3, 2) \geq 19$  (which also implies  $g(3, 3) \geq 38$ ), thus our lower bound is not always optimal.

We would like to conclude with several interesting questions.

**Problem 1.** Is  $g(a, a) = 2g(a - 1, a)$  for all  $a \geq 2$ ?

**Problem 2.** Is  $g(a, b) < 2\binom{a+b}{a}$ ?

**Problem 3.** Is  $g(a, a) = o(2^{2a})$ ?

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